Two stories about love (Distributed Computing and Topology)

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From the book coauthored with Maurice Herlihy and Dmitry Kozlov to be published by Elsevier

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Yet, the characters need to solve some task

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> It is much harder to pursue concurrent activities, such as preparing a ten-course meal with limited pots and pans, all while speaking to a friend on the telephone.

This talk is about

Theory of concurrency

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Using topology

Why

Many models, appear to have little in common besides the common concern with complexity, failures and timing.

Combinatorial topology provides a common framework that unifies these models.

How

Concurrency is challenging because each process, has only a *limited view* of the world (overall state of the computation)

Placing together all these views yields a simplicial complex

Combinatorial Topology

- Discrete approximation of a geometric object
- To study properties invariant under continuous deformations



The stories

• Cheating wives

(A.k.a. muddy children, from knowledge theory)

• Two insecure lovers

(A.k.a. Coordinated attack, from databases and networking)

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- The King of the country announced "There is at least one unfaithful wife" and publicized the following decree

Cheating wives decree

He asks the following question over and over:

can you tell for sure whether or not you are a cuckold?

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Assuming that all of the men are intelligent, honest, and answer simultaneously, what will happen?

Analysis of the puzzle

First operational, then combinatorial

First, suppose that exactly <u>one</u> is cuckold

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- At the first question, exactly one says "yes"
- At the second, all others say "no"

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- At the third, all others say "no"

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Operational analysis (3)

Suppose that exactly k are cuckolds, by induction...

- At the k-th question, exactly k say "yes"
- At the (k+1)-th, all others say "no"

Local states

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- It is represented by a vector: in position i has 0 if man i is known to be clean, and 1 if cuckold
- Secause man *i* does not know its own status, its input vector has \perp in position i

Global inputs

Each possible input configuration is represented as a simplex, linking compatible states for the men



meaning that the men can be in these states together











initial state, except for the one where all are clean

















3 vertexes exposed, where someone knows its status



⊙ 2

0

















All 3 announce "cuckolds"



















No decisions









3 vertexes labeled, "cuckold"



0

2

0

3

























Output complex



Each man should say "yes" or "no" All combinations are possible...

Output complex



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Solving the cheating wives task

Each man decides an output value, on one of its local states

Decisions define a simplicial map from input complex to output complex that respects the task's specification

In this task communication is very limited. More generally, for any task...

Solving any task

In the basic, wait-free model

A task is solvable if and only if there exists a *subdivision* of the input complex and a simplicial map to the output complex that respects the task's specification

Herlihy, Shavit 1993

Wait-free: asynchronous model where any number of processes can crash

Coordination

We often need to ensure that two things happen together or not at all.

For example, a banking system needs to ensure that if an automatic teller dispenses cash, then the corresponding account balance is debited, and viceversa.

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- If both attend, they win, but if only one attends, defeat and humiliation is felt.
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- Communication is be SMS only.

Communication problems

- Normally, it takes a message one hour to arrive.
- However, it is possible that it is gets lost.

The puzzle

Fortunately, on this particular night, all the messages arrive safely.

How long will it take Alice and Bob to coordinate their meeting?
Analysis of the puzzle

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Suppose Alice initiates the communication

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- Although her message was in fact delivered, Alice does not know. She therefore considers it possible that Bob did not receive the message.
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- Showing this, Bob will not show up based solely on Alice's message.

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- Solution Naturally, Bob reacts by sending an acknowledgment back to Alice, which arrives at 2:00
- Will Alice plan to show up?
- Unfortunately, Alice's predicament is similar to Bob's predicament at 1:00, she cannot yet decide to show up

No number of successfully delivered acknowledgments will be enough to ensure that show up safely!

The key insight is that the difficulty is not caused by what actually happens (all messages actually arrive) but by the uncertainty regarding what might have happened.

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- This vertex belongs to two edges (simplexes)



Topology implies impossibility

No number of successfully delivered acknowledgments will be enough to ensure that show up safely, because the complex is subdivided, and remains connected! No number of successfully delivered acknowledgments will be enough to ensure that show up safely!



Because not possible to map a connected input complex into a disconnected output complex



To conclude

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- We have used a variety of topology techniques to analyze concurrency: homology, covering spaces, shellability, decidability, sperner's lemma, etc.

The End Thanks for you attention